

Real-Time Systems: Examples / Case Studies

- Simple Control System
- Sampling Periods
- Quality of the Control vs. Processing Cost
- Protection of Resources in Integrated Systems
- Multimedia / Real-Time Communication
- Anomalies in Asynchronous Systems
 - Example: Advanced Fighter Technology Integration (AFTI) F16
- Priority Inversion

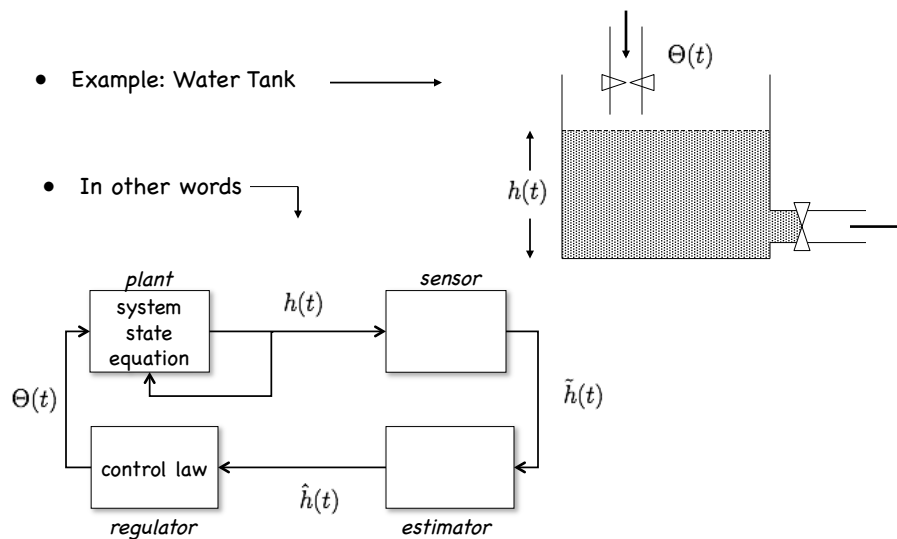
- Real-Time Systems
- Hard and soft deadlines; operational definition

© R. Bettati

Application Areas: Control Systems

- Example: Water Tank

- In other words



© R. Bettati

Control Systems (cont)

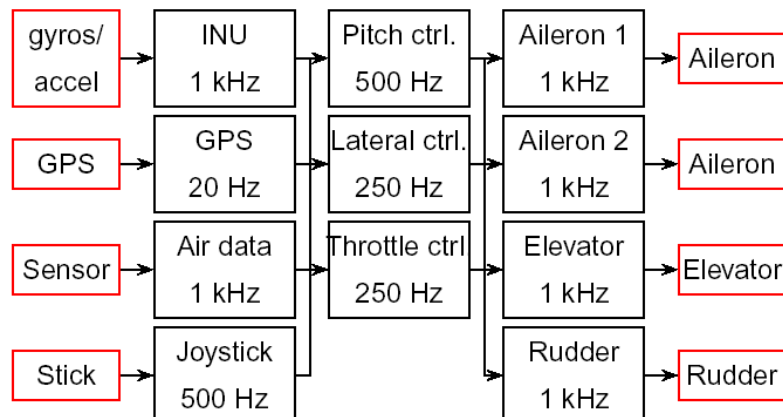
- Control Loop:

```
DO FOREVER
  wait_for_delay
  h      := fluid_height
  theta := valve_position
  r      := table_lookup(h, theta)
  IF     r = left THEN turn_left
  ELSE IF r = right THEN turn_right
  ELSE   do_nothing
ENDDO
```

© R. Bettati

Example: Avionics System

Hard real-time system with multirate behavior

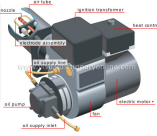


© R. Bettati

Quality of Control vs. Processing Cost Example: Open-Loop Temperature Control

[Simplified from : Setol, Lehoczky, Sha, and Shin, "On Task Schedulability in Real-Time Control Systems",
Proceeding of the 1996 IEEE Real-Time Systems Symposium]

System: Temperature of a unit is controlled by a burner.



Dynamic equation:

$$\dot{x} = -ax + bu$$

x - difference between unit and ambient temperature, $x(0) = 0$
 u - control input (rate of heat)

Control Problem: change temperature of unit to x_d within time t_f ; consume minimum amount of fuel. Allow for a tolerance δ .

$$|x(t_f) - x_d| \leq \delta$$

Performance Index $J(u)$ of control system: measure of total cost of control and accuracy generated in time period $[0, t_f]$ by control u .

Generally:

$$J(u) = S(x(t_f), t_f) + \int_0^{t_f} L(x(t), u(t), t) dt$$

Optimal control $u^*(t)$ with performance index J^* .

© R. Bettati

Open-Loop Temperature Control (cont)

Our case: minimize fuel.

$$\min_u J = \frac{1}{2}p(x(t_f) - x_d)^2 + \frac{1}{2} \int_0^{t_f} u^2(t) dt$$

Resulting **optimal control:**

$$u^*(t) = \frac{x_d p a b e^{at}}{a e^{at_f} + p b^2 \sinh(at_f)}$$

Final State:

$$x^*(t_f) = \frac{x_d p b^2 \sinh(at_f)}{a e^{at_f} + p b^2 \sinh(at_f)}$$

© R. Bettati

Open-Loop Temperature Control (cont)

Discretize control input u :

Sampling period P .

$$\dot{x}^*(t) = -ax^*(t) + bu^*(kP) \quad kP \leq t \leq (k+1)P$$

Performance index for discrete optimal control:

$$J_D^*(P) \doteq S(x^*(t_f), t_f) + \sum_{k=0}^{n-1} \int_{kP}^{(k+1)P} L(x^*(t), u^*(kP), t) dt$$

In our case:

$$J_D^*(P) \doteq \frac{1}{2} px_d \left(\frac{1 - e^{-aP}}{1 + e^{-aP}} \right)$$

Constraints:

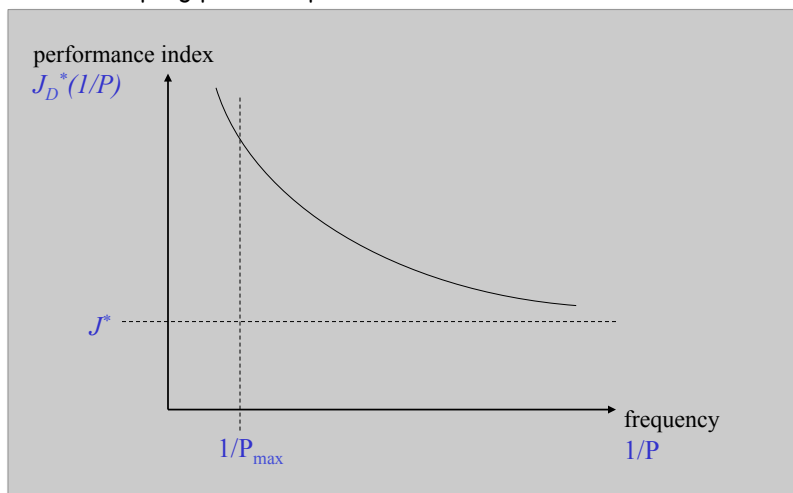
$$|x(t_f) - x_d| \leq \delta$$

$$x_d \left(\frac{1 - e^{-aP}}{1 + e^{-aP}} \right) \leq \delta \Rightarrow P \leq \frac{1}{a} \ln \frac{x_d + \delta}{x_d - \delta}$$

© R. Bettati

Open-Loop Temperature Control (cont)

- Effect of sampling period on performance index.



© R. Bettati

Quality of Control vs. Processing Cost (cont)

Task frequencies must be determined to optimize the performance indices without overloading the available processing capabilities.

Notation: $\Delta J^*(P) := J_D^*(P) - J^*$

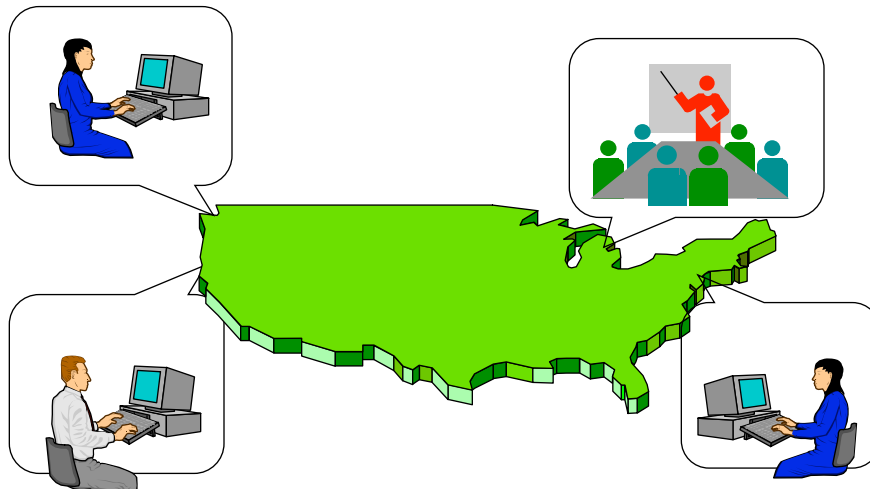
Optimization problem:

Given a set of tasks, T_1, \dots, T_n , with given $\Delta J_i^*(\bullet)$ and execution times C_i , find a set of periods P_i , such that

1. $P_i \leq P_i^{max}$ // Maintain stability
2. Minimize (maximize) $\sum_{i=1}^n \Delta J_i^*(P_i)$ // Optimize total performance index
3. Resource Constraint: $\sum_{i=1}^n \frac{1}{P_i} C_i \leq U$ // CPU capacity

© R. Bettati

Example: Multimedia



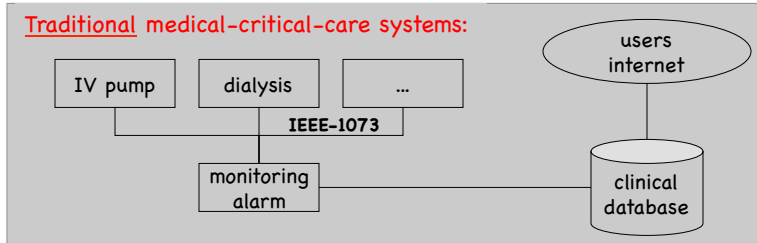
Example: Teleseminars

© R. Bettati

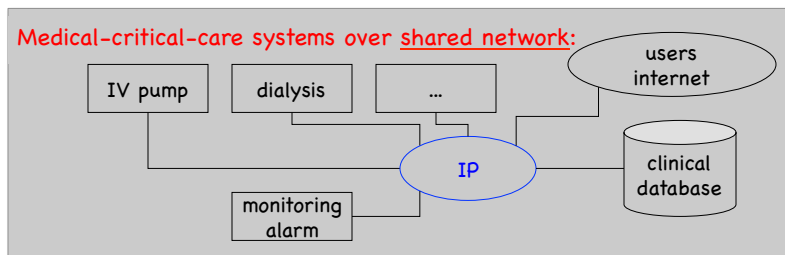
Example: Intensive Care Computing

(Ken Birman, "The Next-Generation Internet: Unsafe at any Speed?", IEEE Computer Aug 2000)

Traditional medical-critical-care systems:

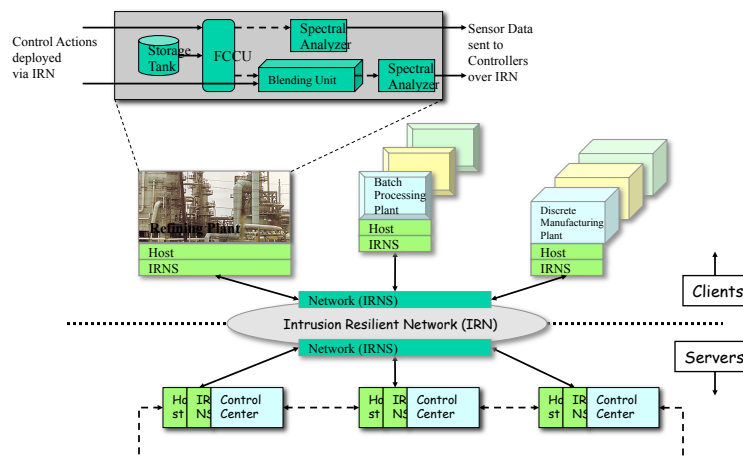


Medical-critical-care systems over shared network:



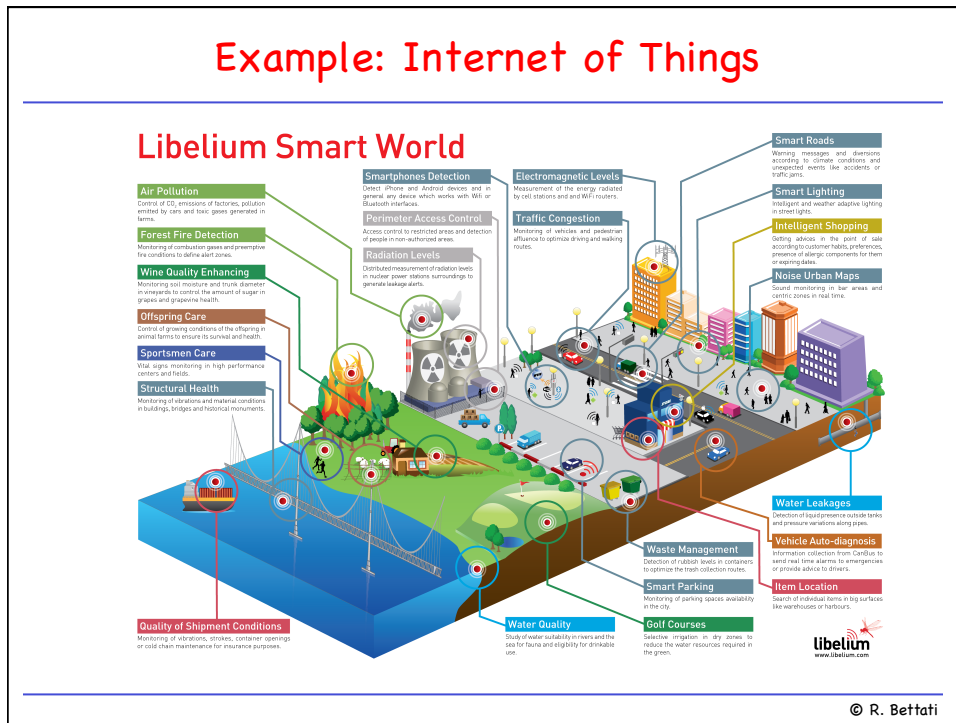
© R. Bettati

Example: Industrial Applications

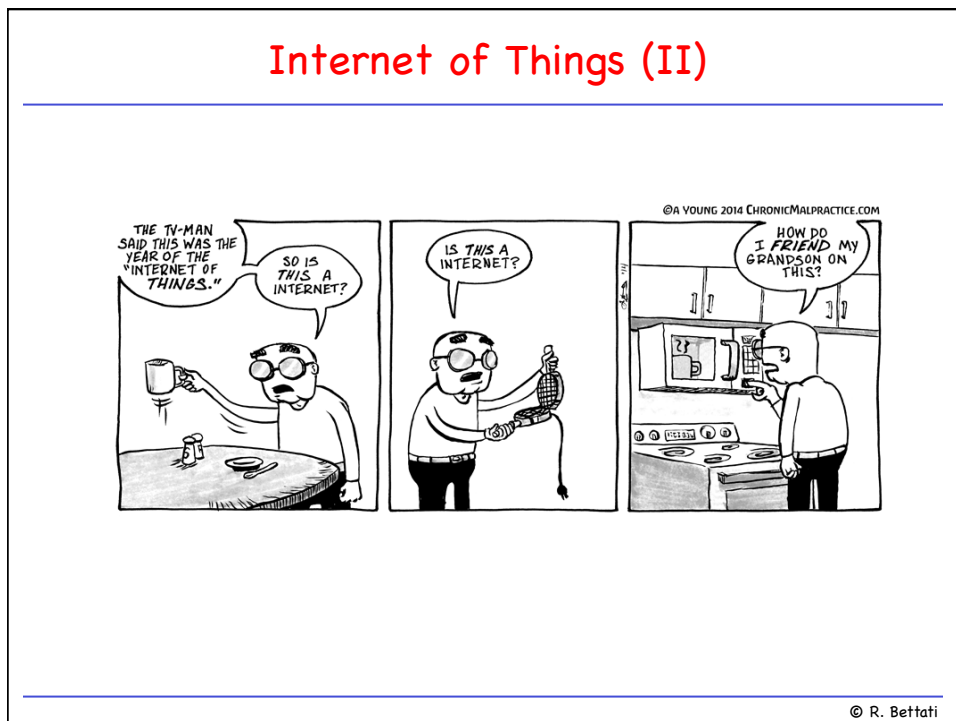


© R. Bettati

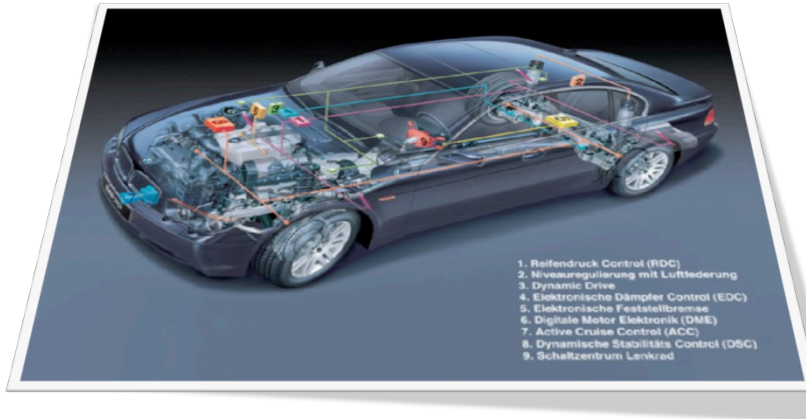
Example: Internet of Things



Internet of Things (II)




Example: Cars as Systems-of-Systems



© R. Bettati

Cars as System of Systems (II)

**Das Elektronik-Konzept
der neuen 7er-Reihe**



Bordnetz

Highlights:
iDrive, Dynamic Drive, Elektrische Feststellbremse,

Bussysteme:	4 Hauptbussystem + Subbusse
Diagnose-Bus-Zugang:	BMW-Fast 115kbaud
Anzahl Steuergeräte:	45 – 75 je nach Ausstattung
Verwendete Prozessoren:	im wesentlichen PPC, HC-12, C167
Betriebssystem:	CAN-BMW Standard-Core / MOST Referenzapplikationen ISIS proprietär
Flash-EPROM:	> 100 Mbyte je nach Ausstattung
Anzahl elektrische Motoren:	Bis zu 150
Leitungslänge Kupfer / LWL:	ca. 2,5 km / ca.50 m (ISIS 43m + MOST 7m)
Signale auf dem Bus:	ca. 2500
Anzahl Lieferanten:	mehr als 15 mit jeweils mehreren Standorten
Powermanagement:	Aufstart-/Abschaltverhalten und Ruhestrom

→ Herausforderung Integration

BMW AG Axel Deicke EE-I
FH Deggendorf AK Mechatronik
Seite 34

© R. Bettati

Cars as SoS (III)

Das Elektronik-Konzept der neuen 7er-Reihe

Bordnetz

System-CAN Ausstattung	Peripherie-CAN Ausstattung	Powertrain-CAN Antrieb/Fahrwerk	byteflight Sicherheit	MOST Kommunikation
<ul style="list-style-type: none"> EMV-Festigkeit, Fehlertoleranz, Ereignisgesteuert 100 kB/s 	Wie System K-CAN + Entkopplung des Gefahrenbereichs Crash, „Angrifer“	<ul style="list-style-type: none"> Regelvorgänge Zyklisch + Ereignisgest. 500 kB/s 	<ul style="list-style-type: none"> System ISIS Sicherheit Teilnetzfähig 10 MB/s 	<ul style="list-style-type: none"> Kontinuierliche Signale Automotivstandard ca.22 MB/s
<p>Baumstruktur Kupferkabel</p>	<p>Baumstruktur Kupferkabel</p>	<p>Baumstruktur Kupferkabel</p>	<p>Sternstruktur Lichtwellenleiter</p>	<p>Ringstruktur Lichtwellenleiter</p>

- **Integrierter Systementwurf:** Homogenität über alle Busse
- **Powermodul** stellt Startfähigkeit selbst in Fehlerfällen sicher, Abschaltung von Verbrauchern
- **Standard Core:** Vereinheitlichung der STG-Basisfunktionalität.
- **Gateway:** Firewall zwischen den Teilsystemen

BMW AG Axel Daicke EE-I
FH Deggendorf AK Mechatronik
Seite 35

© R. Bettati

Cars as SoS (IV)

Das Elektronik-Konzept der neuen 7er-Reihe

Bordnetz

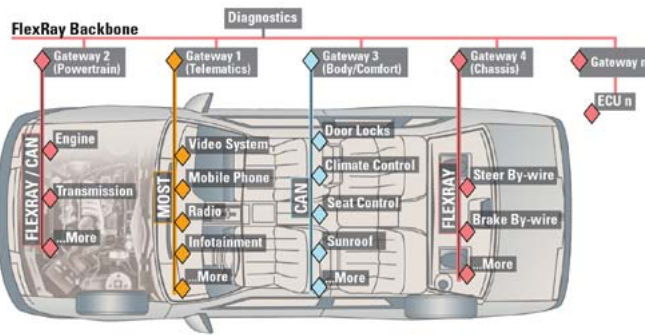
K-CAN System
MOST
K-CAN Peripherie
byteflight
PT-CAN

BMW AG Axel Daicke EE-I
FH Deggendorf AK Mechatronik
Seite 36

© R. Bettati

Cars as SoS

Example of a Backbone Architecture with FlexRay



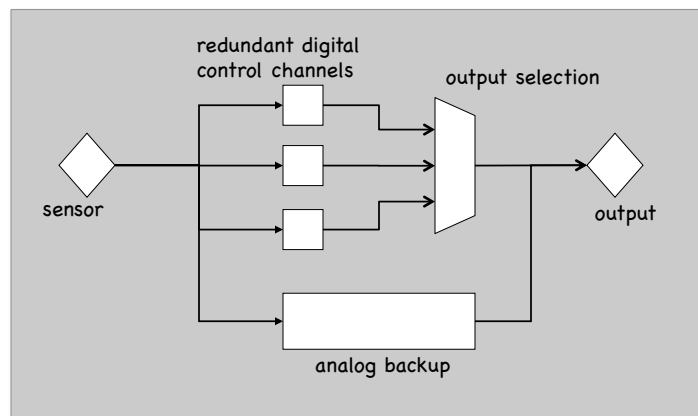
(www.autofieldguide.com)

© R. Bettati

Example: Asynchronous Design of Digital Flight Control Systems

(J. Rushby, SRI-CSL-93-07, Nov. 1993)

- Advanced Fighter Technology Integration (AFTI) F-16 DFCS:



© R. Bettati

Asynchronous Design of Digital Flight Control Systems

“... The **asynchronous design** of the [AFTI-F16] DFCS introduced a **random, unpredictable** characteristic into the system. The system became **untestable** in that testing for each of the possible time relationships between the computers was impossible. This random time relationship was a major contributor to the **flight test anomalies**. Adversely affecting testability and having only postulated benefits, asynchronous operation of the DFCS demonstrated the **need to avoid random, unpredictable, and uncompensated design characteristics.**”

D. Mackall, flight-test engineer AFTI-F16 flight tests

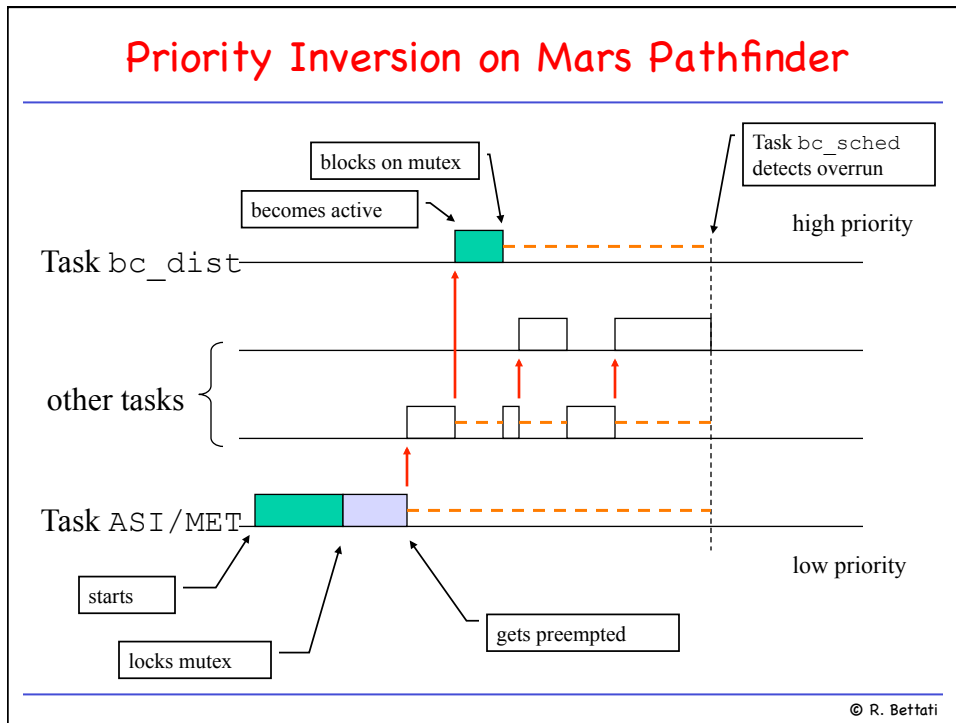
© R. Bettati

Example: Mars Pathfinder Incident



- Landing on July 4, 1997
- “experiences software glitches”
- Pathfinder experiences repeated RESETS after starting gathering of meteorological data.
- RESETS generated by watchdog process.
- Timing overruns caused by priority inversion.
- Resources:
http://research.microsoft.com/en-us/um/people/mbj/Mars_Pathfinder/Mars_Pathfinder.html

© R. Bettati



Real-Time vs. Non-Real-Time Systems

Q: What distinguishes RT systems from non-RT systems?

A: Timing Constraints!

© R. Bettati

Players in Real-Time Systems

Jobs and Processors:

- **Job:** Unit of work executed by the system
 - **Processor:** Jobs require resource to execute (CPU, disk, network link)
- (We don't distinguish between types of processors!)

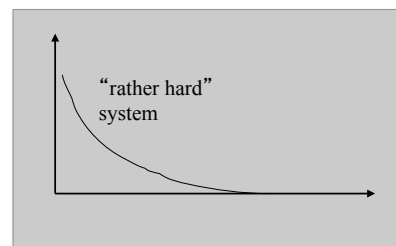
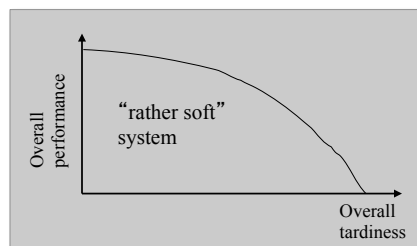
Timing constraints:

- **Release Time:** time when job becomes available for execution
- **Deadline:** time when execution must be completed
- **Relative Deadline:** maximum response time

© R. Bettati

Hard vs. Soft Deadlines

- **Hard Deadline:** Late result may be a fatal flaw, of little use, or cause disastrous consequences
- **Soft Deadline:** Timely completion desirable. Late results useful to some degree.
- **Quantitative measure:** Overall system performance as function of tardiness of jobs.



Operational Definition: A job has a **hard** deadline whenever the system designer must **prove** that the job **never** misses its deadline.

© R. Bettati

Hard Real-Time Systems

Definition: A real-time system is **hard-real-time** when a large portion of the deadlines is hard.

- Examples:
 - Embedded systems
 - Recovery procedures in high-availability systems
 - many others ...

© R. Bettati

Hard Real-Time Systems

Q: Does real-time mean fast ?

© R. Bettati

Hard Real-Time Systems

Q: Why not use commercial (general purpose) OSs?

A: Verification, Certification

© R. Bettati

Hard Real-Time Systems

Q: Why do we need to meet deadlines 100% of the time?

A:

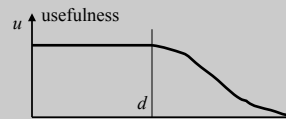
- Validation of probabilistic timing requirements.
- Assessment of **compound effect** of missed deadlines with other factors.

© R. Bettati

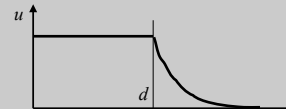
Soft Real-Time Systems

Definition: A real-time system is **soft-real-time** when jobs have soft deadlines.

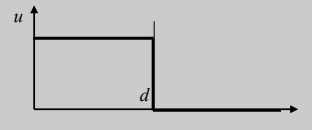
Non-stringent timing requirements
on-line transaction system
telephone switches



More stringent timing requirements
Stock price quotation system



Stringent timing requirements
Multimedia



Requirements often specified in **probabilistic** terms; **validation** is often done by **simulation**, **trial use**.