NP-Completeness

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[partially based on slides by Jennifer Welch]
Definition of NP-Complete

L is NP-complete if and only if

(1) L is in NP and

(2) for all L' in NP, L' \leq_p L.

In other words, L is at least as hard as every language in NP.
Theorem  Suppose L is NP-complete.

(a) If there is a poly time algorithm for L, then $P = NP$.

(b) If there is no poly time algorithm for L, then there is no poly time algorithm for any NP-complete language.
Proving NP-Completeness

(a) Use a direct approach and prove that
   (1) L is in NP
   (2) every other language in NP is polynomially reducible to L

(b) Find an NP-complete problem and use reduction.

Approach (a) is for larger-than-life people, (b) is for mere mortals.
Proving NP-Completeness by Reduction

To show $L$ is NP-complete:

1. Show $L$ is in NP.

2.a) Choose an appropriate known NP-complete language $L'$.

2.b) Show $L' \leq_p L$.

This works, since every language $L''$ in NP is polynomially reducible to $L'$, and $L' \leq_p L$. By transitivity, $L'' \leq_p L$. 
SAT
First NP-Complete Problem

How do we get started? Need to show via brute force that some problem is NP-complete.

- Logic problem "satisfiability" (or SAT).

- Given a boolean expression (collection of boolean variables connected with ANDs and ORs), is it satisfiable, i.e., is there a way to assign truth values to the variables so that the expression evaluates to TRUE?
Conjunctive Normal Form (CNF)

**Boolean variable:** Indeterminate with values T or F. Example: \( x, y \)

**Literal:** Variable or negation of a variable. Example: \( x, \neg x \)

**Clause:** Disjunction (OR) of several literals. Example: \( x \lor \neg y \lor z \lor w \)

**CNF formula:** Conjunction (AND) of several clauses. Example: \( (x \lor y) \land (z \lor \neg w \lor \neg x) \)
Satisfiable CNF Formula

• Is \((x \lor \neg y)\) satisfiable?
  - yes: set \(x = T\) and \(y = F\) to get overall T
• Is \(x \land \neg x\) satisfiable?
  - no: both \(x = T\) and \(x = F\) result in overall F
• Is \((x \lor y) \land (z \lor w \lor x)\) satisfiable?
  - yes: \(x = T, y = T, z = F, w = T\) result in overall T
• If formula has \(n\) variables, then there are \(2^n\) different truth assignments.
Definition of SAT

SAT = all (and only) strings that encode satisfiable CNF formulas.
SAT is NP-Complete

- Cook's Theorem: SAT is NP-complete.
- Proof ideas:
  - (1) SAT is in NP: Given a candidate solution (a truth assignment) for a CNF formula, verify in polynomial time (by plugging in the truth values and evaluating the expression) whether it satisfies the formula (makes it true).
SAT is NP-Complete

- How to show that every language in NP is polynomially reducible to SAT?

- Key idea: the common thread among all the languages in NP is that each one is solved by some nondeterministic Turing machine (a formal model of computation) in polynomial time.

- Given a description of a poly time TM, construct in poly time, a CNF formula that simulates the computation of the TM.