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```
Input: G = (V,E)
for each node u in V do
    mark u as unvisited
od;
for each unvisited node u do
   recursiveDFS(u);
od;
```

```
recursiveDFS(u):

mark u as visited;

for each unvisited neighbor v of u do

recursiveDFS(v)

od
```

Purpose of this loop: Create DFS forest. Graph can be directed or undirected.

DFS Forest

By keeping track of parents, we want to construct a forest resulting from the DFS traversal.

```
Input: G = (V,E)
for each node u in V do
     parent[u] = nil;
     mark u as unvisited
od;
for each unvisited node u do
     parent[u] = u;
     recursiveDFS(u);
od;
```

```
recursiveDFS(u):
   mark u as visited;
   for each unvisited neighbor v of u do
      parent[v] = u; recursiveDFS(v)
   od
```

Refining DFS

Let us keep track of some interesting information for each node. We will timestamp the steps and record the

- · discovery time, when the recursive call starts
- · finish time, when its recursive call ends

```
Input: G = (V,E)
for each node u in V do
     parent[u] = nil;
     mark u as unvisited
od;
time = 0;
for each unvisited node u do
     parent[u] = u;
     recursiveDFS(u);
od;
```

```
recursiveDFS(u):
    mark u as visited;
    disc[u] = ++time;
    for each unvisited neighbor v of u do
        parent[v] = u; recursiveDFS(v)
    od;
    fin[u] = ++time;
```

Running Time of DFS

The first for-loop for initialization takes O(V) time.

The second for-loop in non-recursive wrapper considers each node, so O(V) iterations.

One recursive call is made for each node. In a recursive call for the node u, all its neighbors are checked; so the total time in all recursive calls is O(E).

Total time is O(V+E).

Nested Intervals

Define [disc[v], fin[v]] to be the interval for node v.

Claim: For any two nodes, either one interval precedes the other or one is enclosed in the other

Indeed, the recursive calls are nested.

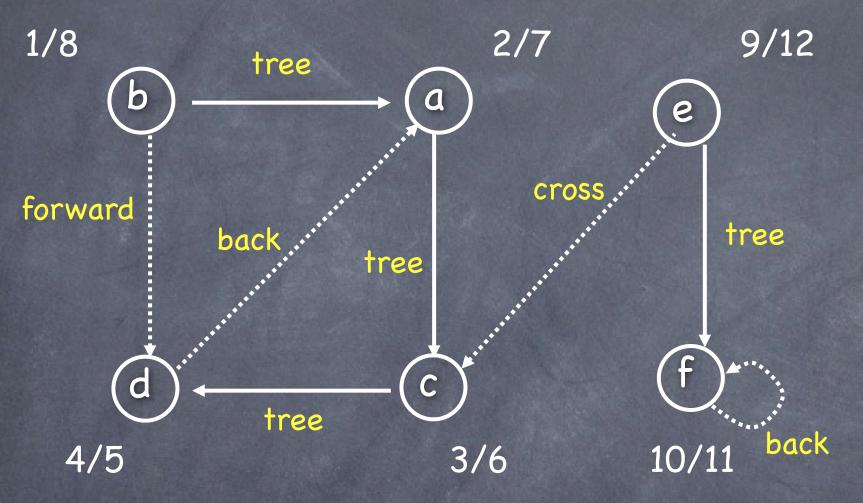
Corollary: v is a descendant of u in the DFS forest iff the interval of v is inside the interval of u.

Classifying Edges

Consider edge (u,v) in a directed graph G = (V,E) with respect to its DFS forest

- otree edge: v is a child of u
- Oback edge: v is an ancestor of u
- oforward edge: v is a descendant of u but not a child
- ocross edge: none of the above

Example of Classifying Edges



in DFS forest

····· not in DFS forest

a/b disc./finish. time

tree edge: v child of u back edge: v ancestor of u forward edge: v descendant of u, but not child cross edge: none of the above