Cosimulation II

Cosimulation Approaches

How to cosimulate?

- *How to simulate hardware components of a mixed hardware-software system within a unified environment?*
  - This includes simulation of the hardware module, the processor, and the software that the processor executes.
- *How to simulate hardware and software at same time?*
- What are various challenges?
  - Software runs faster than hardware simulator. How to run the system simulation fast keeping the above synchronized?
  - Slow models provide detailed and accurate results than fast models. How to balance these effects?
  - Use of different platforms for simulations.
Some basic approaches

• Detailed Processor Model:
  – processor components (memory, datapath, bus, instruction decoder etc) are discrete event models as they execute the embedded software.
  – Interaction between processor and other components is captured using native event-driven simulation capability of hardware simulator.
  – Gate level simulation is extremely slow (~tens of clock cycles/sec), behavioral model is ~hundred times faster. Most accurate and simple model

Some basic approaches

• Bus Model (Cycle based simulator):
  – Discrete-event shells that only simulate activities of bus interface without executing the software associated with the processor. Useful for low level interactions such as bus and memory interaction.
  – Software executed on ISA model and provide timing information in clock cycles for given sequence of instructions between pairs of IO operation.
  – Less accurate but faster simulation model.
Some basic approaches

• Instruction Set Architecture Model:
  – ISA can be simulated efficiently by a C program. C program is an interpreter for the embedded software.
  – No hardware mode. Software executed on ISA model. Execution on ISA model provides timing (clock) details of the cosimulation.
  – Can be more efficient than detailed processor modeling because internals of the processor do not suffer the expense of discrete-event scheduling.

• Compiled Model:
  – very fast processor models are achievable in principle by translating the executable embedded software specification into native code for processor doing simulation. (Ex: Code for programmable DSP can be translated into Sparc assembly code for execution on a workstation)
  – No hardware, software execution provides timing details on interface to cosimulation.
  – Fastest alternative, accuracy depends on interface information.
Some basic approaches

- Hardware Model:
  - If processor exists in hardware form, the physical hardware can often be used to model the processor in simulation. Alternatively, processor could be modeled using FPGA prototype. (say using Quickturn)
  - Advantage: simulation speed
  - Disadvantage: Physical processor available.

A New Approach

- This is a combined HW/SW approach. The host is responsible of having OS, some applications and might have superset simulating environment (RSIM, SIMICS, SIMOID).
- Use of fast backplane (PCI) for communication. Real processor or processor core in FPGA as hardware model, and ASIC/FPGA for interface and interconnection for hardware modeler.
- Good for fast complex architecture simulations including multiprocessor.
Domain coupling

- In four out of six approaches, we have host that run software and required to interact with hardware model or simulator.
- Difficulties:
  - Providing timing information across the boundaries
  - Coupling two domains with proper synchronization

Migration across cosimulation

- Consider the system simulation at different levels of abstraction throughout the design process:
  - In the beginning of design process, hardware synthesis is not available. Hence use functional model to study the interaction between HW and SW.
  - As design progress with more implementations, replace functional model of hardware by netlist level.
  - Once detail operation of hardware is verified, swap back the high level description of HW design to gain simulation speed.
- The cosimulation environment should have this migration support across the levels of abstraction.
- Off-the-shelf Components: design is not a part of the current design process. Functional model is enough, no need to know internal details.
Master slave cosimulation

- One master simulator and one or more slave simulators: slave is invoked from master by procedure call.
- The language must have provision for interface with different language.
- Difficulties:
  - No concurrent simulation possible
  - C procedures are reorganized as C functions to accommodate calls

Distributed cosimulation

- Software bus transfers data between simulators using a procedure calls based on some protocol.
- Implementation of System Bus is based on system facilities (Unix IPC or socket). It is only a component of the simulation tool.
- Allows concurrency between simulators.
Synchronization and Time in cosimulation

- In case of a single simulator (say Verilog) there is no problem for timing as single event queue is managed for simulation.
- If there are several simulators and software programs in the domain:
  - hardware and software domain are using a handshaking protocol to keep their time (clock) synchronized. Signals (events) transferred from one side to the other should have attached a time stamp.
  - It is possible to use a loosely coupled strategy which allows the two domain to proceed more independently. If a signal is received with a time stamp lower than the current clock in the respective domain, the respective simulator have to be back up.

Aspects of cosimulation

- A frame work of cosimulation consists of variety of components, levels of abstractions and different models.

[Diagram showing the framework and levels of abstractions]

Functional Model

Instruction-set model

Detail arch. model

Analog model

Netlist-level
Heterogeneous Environment: Ptolemy

- Ptolemy supports different design styles encapsulated in objects called **Domain**.
  - Domain realizes a computational model appropriate for a particular subsystem. Ex: SDF, Dynamic Dataflow(DDF), Discrete Event(DE) and Digital hardware modeling environment (Thor).
- Domain consists of a set of Blocks and Schedulers that confirm to a common computational model.

**Block**: setup, go(), wrapup().

**Geodesic**: initialize(), numInit(), setSourcePort(), setDestPort().

**Particle**: readType(), print(), operator <<(), clone().

**PortHole**: initialize(), receiveData(), sendData().

*Block objects send and receive data encapsulated in particles to outside world through PortHoles.*

*Buffering: by Geodesic*

*Garbage Collection: Plasma*

Heterogeneous cosimulation: Ptolemy

- Any model can be used at the top of hierarchy. Within each level of hierarchy, it is possible to have Blocks containing foreign domain.
- Hierarchy heterogeneity is quite different from the concept of a simulation backplane (that imposes a top-level models of computation through which all subsystem interact).
- Active objects in a Domain are **Stars**. They perform computation and communications with other objects through PortHoles.
Heterogeneous cosimulation: Ptolemy

- Domain XXX can contain foreign Domain YYY in it. Call it XXXWormhole
- To the stars of Domain XXX, the XXXWormhole appears as another star. Though, inside the Wormhole, it is entirely another domain.

Cosimulation: Ptolemy

- EventHorizon provides an interface between the external and internal Domains (Wormholes).
- There exists an EventHorizon for all Domains. Thus each domain needs only an interface to this EventHorizon.